

# Tensors: Entanglement, Geometry, and Combinatorics

Benjamin Lovitz

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Department of Applied Math and Institute for Quantum Computing

University of Waterloo

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# Tensors and their decompositions

- Let  $X, Y, Z$  be  $\mathbb{F}$ -vector spaces. A **tensor** is an element  $T \in X \otimes Y \otimes Z$
- A **decomposition** of  $T$  is an expression of  $T$  as a sum of product tensors

$$T = \sum_{a=1}^n x_a \otimes y_a \otimes z_a . \quad (1)$$

- $\text{rank}(T)$  = minimum number of product tensors needed to decompose  $T$ .
- An expression (1) is the **unique rank decomposition** of  $T$  if  $\text{rank}(T)=n$  and any other rank decomposition of  $T$  is the same up to re-ordering summands.

# Outline



## We can view tensors in three different ways:

Chapter 7



- Algebraic statistics: View a **probability vector** for random variables  $X, Y, Z$  as a tensor  $T$ .

Chapter 7

- Algebraic complexity theory: View a **bilinear map**  $X \times Y \rightarrow Z^*$ , e.g. matrix multiplication, as a tensor  $T$ .

Chapters 3-6

- Quantum information: View a (pure) **quantum state** shared by three parties  $X, Y, Z$  as a tensor  $T$ .

# Chapter 7: Algebraic Statistics

Algebraic statistics: View a **probability vector** for random variables  $X, Y, Z$  as a tensor  $T$ .

- Decomposition of  $T$

$$T = \Pr(X, Y, Z) = \sum_l \Pr(l) \Pr(X|l) \otimes \Pr(Y|l) \otimes \Pr(Z|l)$$

- ← Choice of latent random variable  $L$  such that

$$\Pr(X, Y, Z|l) = \Pr(X|l) \otimes \Pr(Y|l) \otimes \Pr(Z|l)$$

$X, Y, Z$  independent conditioned on each outcome  $l$  for  $L$

- Unique rank decomposition of  $T$  → Unique choice of  $L$  with this property.

[Lovitz-Petrov 21]:

- “Splitting theorem” for sets of product tensors.
- Corollary: A generalization of Kruskal’s theorem.

# Generalization of Kruskal's theorem

Generalization of Kruskal's theorem [Lovitz-Petrov 21]: A sufficient condition for

$$T = \sum_{a=1}^n x_a \otimes y_a \otimes z_a$$

to be the unique rank decomposition of T.

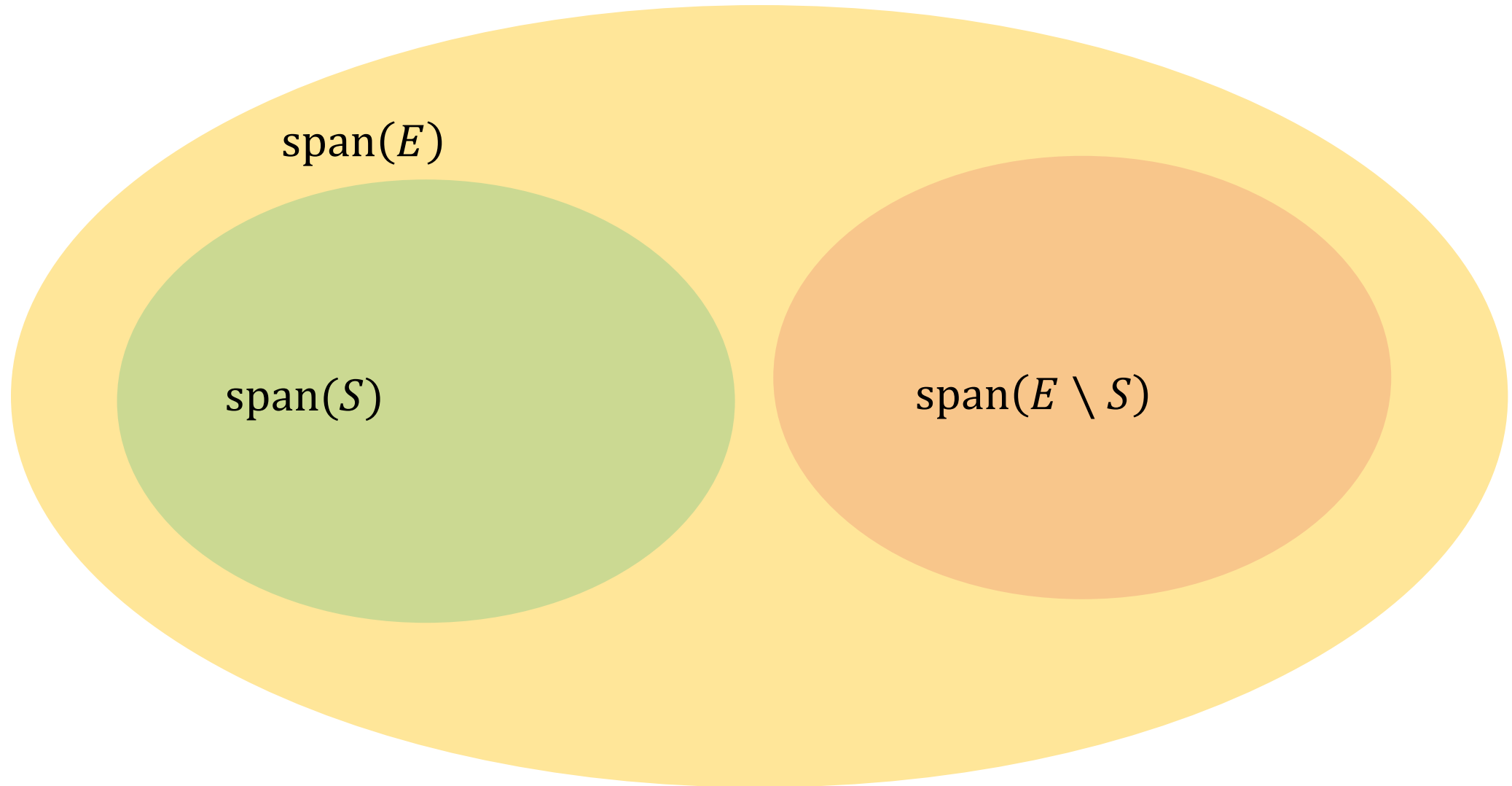
This is a corollary to our splitting theorem...

# Splitting

Definition: A set of vectors  $E = \{v_1, \dots, v_n\}$  **splits** if there exists a non-trivial subset  $S \subseteq E$  such that

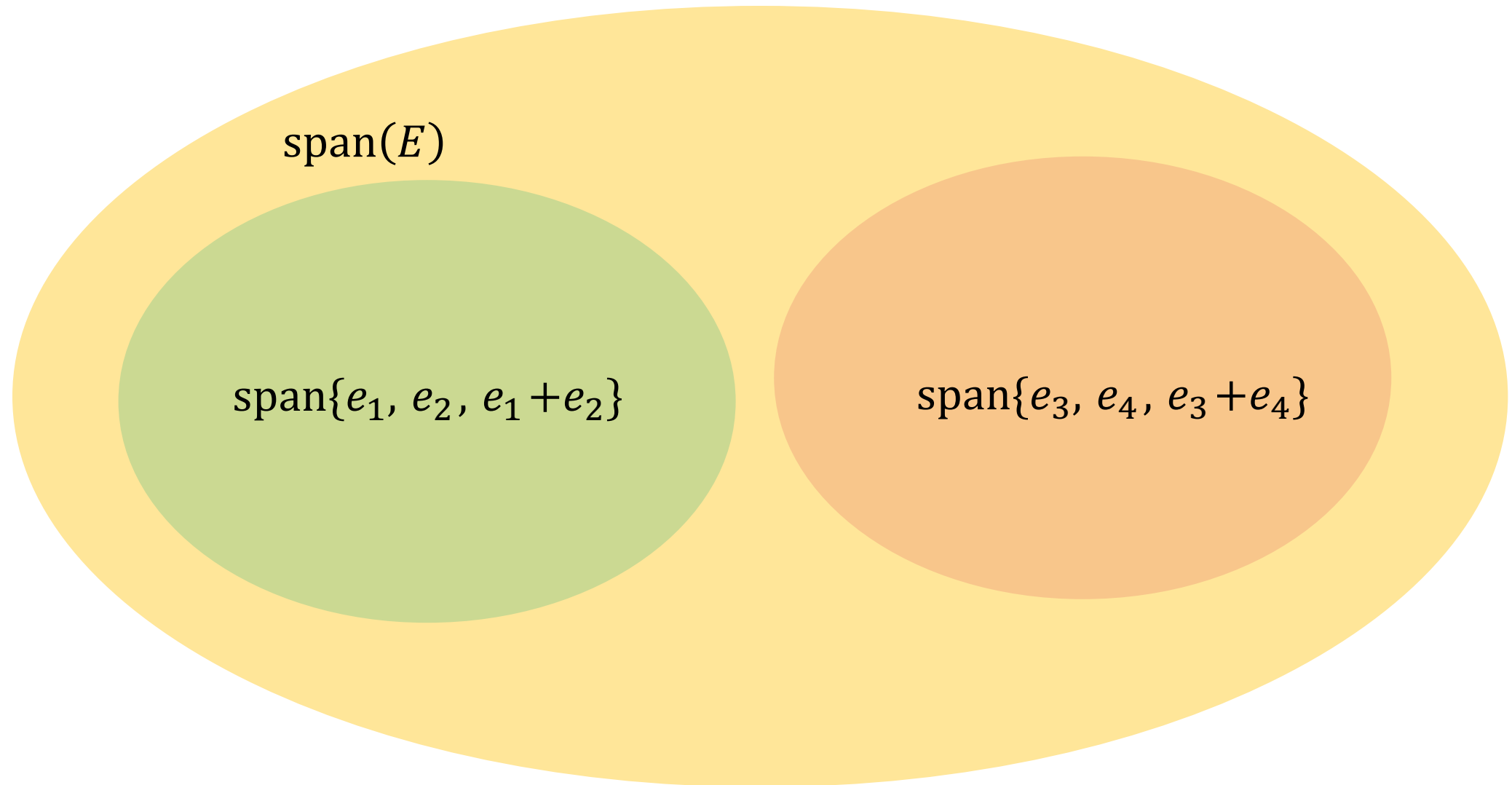
$$\text{span}(S) \cap \text{span}(E \setminus S) = \{0\} \quad (2)$$

$E$  **splits** if there exists  $S \subseteq E$  such that  $\text{span}(S) \cap \text{span}(E \setminus S) = \{0\}$



$$E = \{e_1, e_2, e_1 + e_2, e_3, e_4, e_3 + e_4\}$$

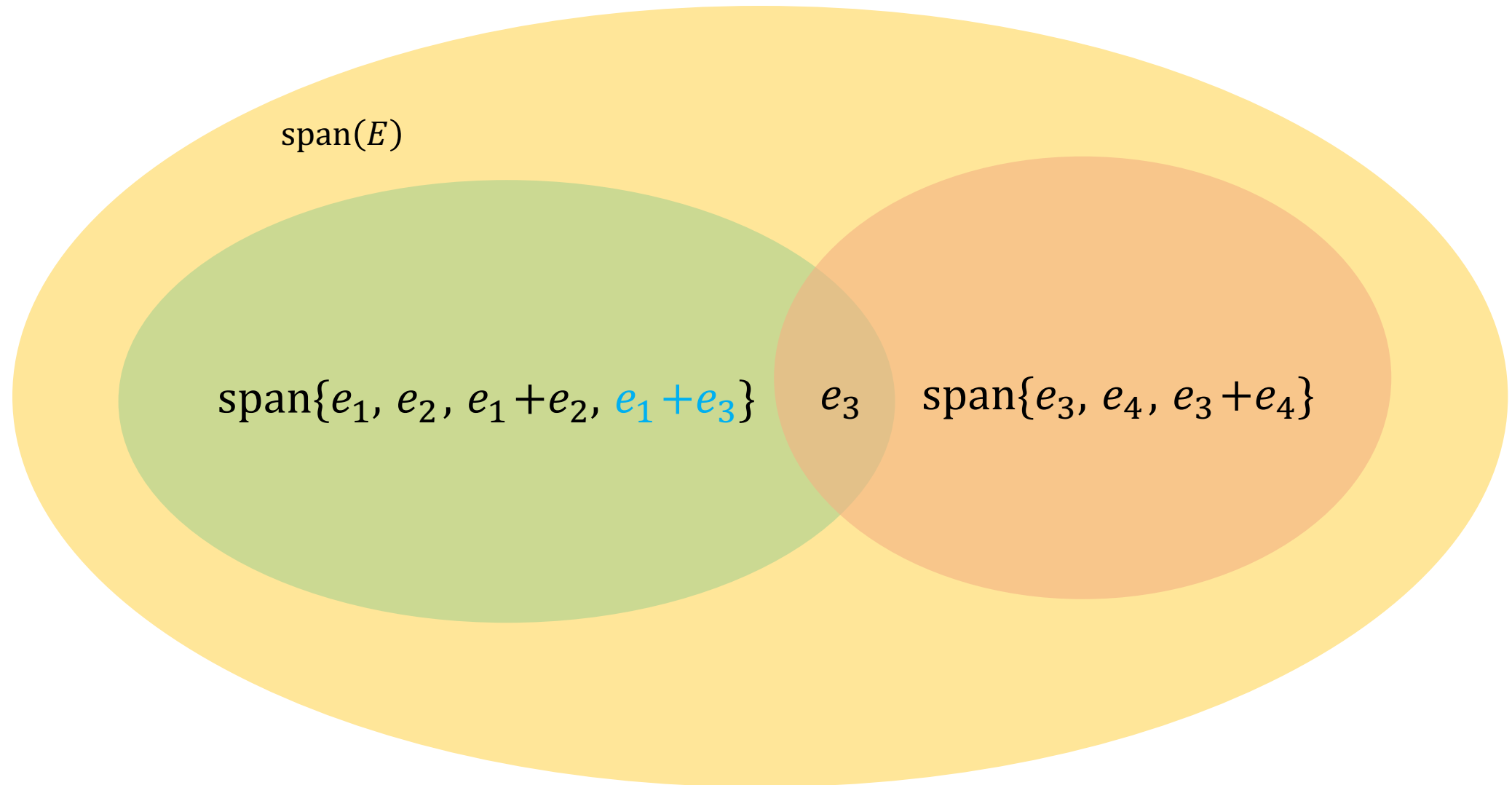
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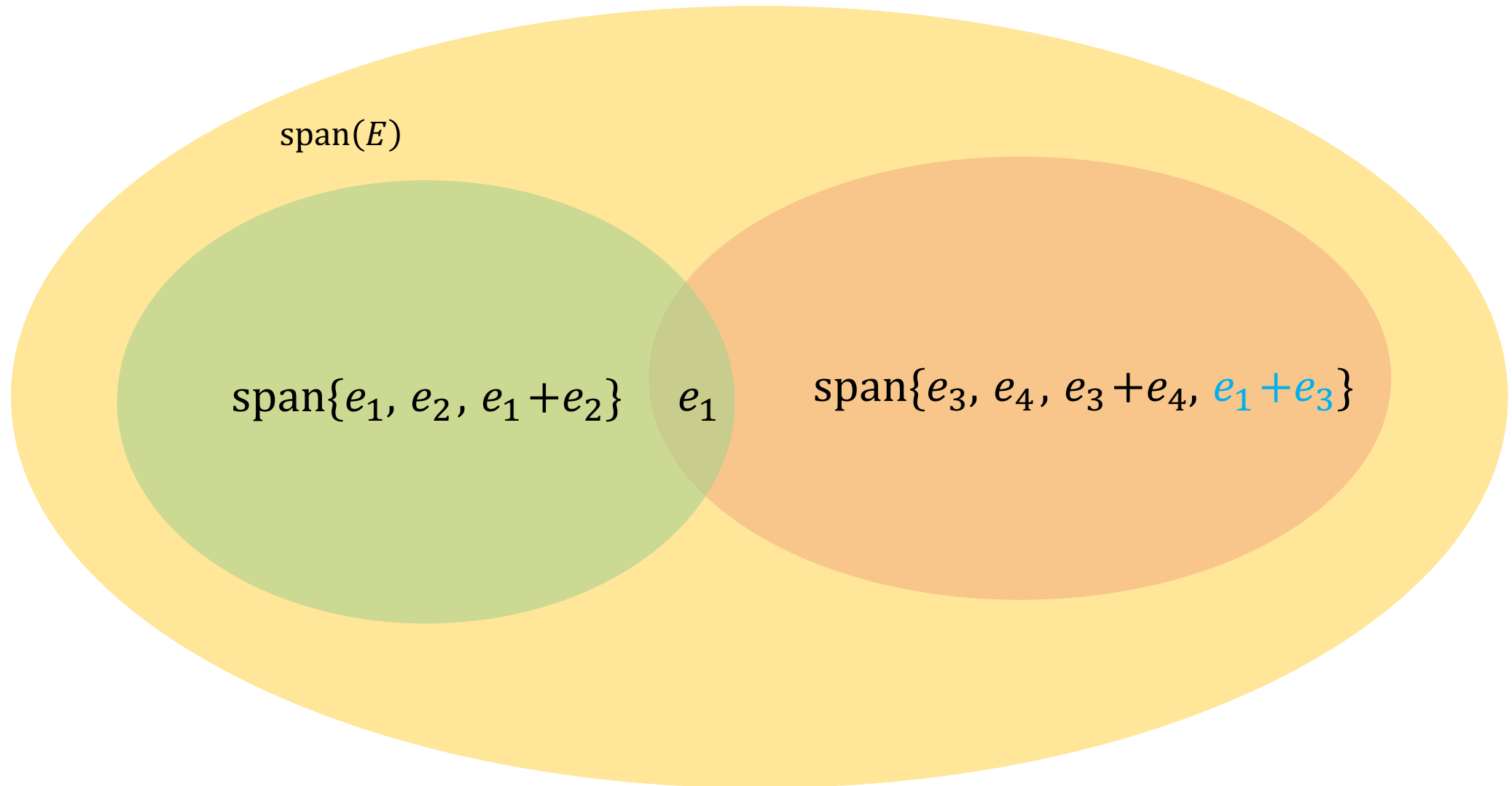
$E$  **splits** if there exists  $S \subseteq E$  such that

$$E = \{e_1, e_2, e_1 + e_2, e_3, e_4, e_3 + e_4\} \cup \{e_1 + e_3\} \quad \text{span}(S) \cap \text{span}(E \setminus S) = \{0\}$$



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# Splitting theorem

$E$  **splits** if there exists  $S \subseteq E$  such that  $\text{span}(S) \cap \text{span}(E \setminus S) = \{0\}$

Splitting theorem [Lovitz-Petrov]: Let  $E = \{x_a \otimes y_a : a \in [n]\}$ .

If

$$\text{dimspan}(E) \leq d_x^{[n]} + d_y^{[n]} - 2$$

then  $E$  splits.


$$d_x^{[n]} = \text{dimspan}\{x_1, \dots, x_n\}$$

(Our generalization of Kruskal's theorem is a corollary to this)

# Splitting theorem

$E$  **splits** if there exists  $S \subseteq E$  such that  $\text{span}(S) \cap \text{span}(E \setminus S) = \{0\}$

Splitting theorem [Lovitz-Petrov]: Let  $E = \{x_a \otimes y_a \otimes z_a : a \in [n]\}$ .

If

$$\text{dimspan}(E) \leq d_x^{[n]} + d_y^{[n]} + d_z^{[n]} - 3$$


then  $E$  splits.


$$d_x^{[n]} = \text{dimspan}\{x_1, \dots, x_n\}$$

# Outline



## We can view tensors in three different ways:

- Chapter 7 • Algebraic statistics: View a **probability vector** for random variables  $X, Y, Z$  as a tensor  $T$ .
- Chapter 7  • Algebraic complexity theory: View a **bilinear map**  $X \times Y \rightarrow Z^*$ , e.g. matrix multiplication, as a tensor  $T$ .
- Chapters 3-6 • Quantum information: View a (pure) **quantum state** shared by three parties  $X, Y, Z$  as a tensor  $T$ .

# Chapter 7: Algebraic Complexity Theory

Algebraic complexity theory: View a **bilinear map**  $X \times Y \rightarrow Z^*$ , e.g. matrix multiplication, as a tensor  $T$ .

- $\text{rank}(T) \propto$  Multiplicative complexity of applying  $T$


[Lovitz-Petrov 21]:

- Corollary to splitting theorem: New lower bounds on tensor rank.

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# Chapters 3-6: Quantum Information

Quantum information: View a (pure) **quantum state** shared by three parties  $X, Y, Z$  as a tensor  $T$ .

- $\text{rank}(T)$  = “GHZ cost” of  $T$  = Smallest size of GHZ state needed to obtain  $T$  by SLOCC.
- Stabilizer rank of  $T$  =  $\chi(T)$  = Computational cost of classically simulating Clifford circuits applied to  $T$  under stabilizer formalism.



# Chapter 6: On Decomposable Correlation Matrices

Question: Given unit vectors  $v_1, \dots, v_n \in \mathbb{C}^d$ , does there exist a positive integer  $m$  and an isometry  $U: \mathbb{C}^d \rightarrow (\mathbb{C}^2)^{\otimes m}$  for which  $Uv_a \in (\mathbb{C}^2)^{\otimes m}$  is a product tensor for all  $a \in [n]$ ?

If “yes”: Protocol using  $v_1, \dots, v_n$  could be converted into a protocol that uses tensor products of qubits.

If “no”: The set  $\{v_1, \dots, v_n\}$  is “absolutely entangled.”

Corollary to splitting theorem: If  $\{v_1, \dots, v_n\}$  does not split, then we can assume  $m \leq n - 2$

# Chapter 6: On Decomposable Correlation Matrices

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Theorem [Lovitz, 2019]:

$n$	1	2	3	4	$\geq 5$
Does such a $U$ always exist?	Yes	Yes	Yes	?	No

# Chapter 3: Generic local state discrimination with pre-shared entanglement

- Question: Given a pure state  $T \in X \otimes Y \otimes Z$ , how useful is it for distributed computation?

- [Lovitz-Johnston 21]: For local unambiguous state discrimination,

$$\dim((GL_X \times GL_Y \times GL_Z)T)$$

is the number of generic pure states in  $X \otimes Y \otimes Z$  that can be unambiguously discriminated with pre-shared entanglement  $T$ .

- Corollary: Generic pure states  $T$  maximize this number.

# Chapter 5: New techniques for bounding stabilizer rank

A state  $S \in (\mathbb{C}^2)^{\otimes m}$  is a **stabilizer state** if  $S = Ue_0^{\otimes m}$  for some Clifford circuit  $U$ .

Recall: A **Clifford circuit** is a circuit composed of Clifford gates

$$H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$$

$$S = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$$

$$CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$$

# Stabilizer rank

- The **stabilizer rank** of a state

$$T \in (\mathbb{C}^2)^{\otimes m},$$

denoted  $\chi(T)$ , is the minimum number  $r$  for which

$$T = \sum_{i=1}^r c_i S_i$$

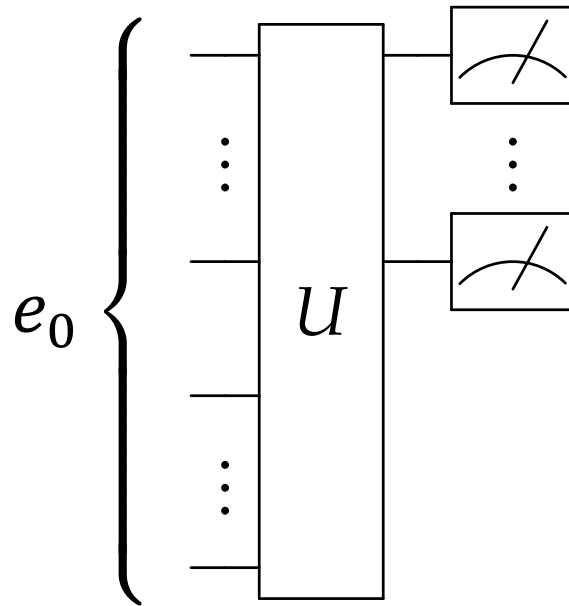
for some  $c_i \in \mathbb{C}$  and  $S_i$  stabilizer states.

- The  **$\delta$ -approximate stabilizer rank** of  $T$  is

$$\chi_\delta(T) = \min\{\chi(R) : \|T - R\| \leq \delta\}.$$

# Motivation: Classical simulation

Question: Given a quantum circuit



Can it be simulated efficiently on a classical computer?

# Motivation: Classical simulation

Partial answer: Write  $U$  as a sequence of Clifford+T gates (universal gate set).

Clifford gates:  $H = \frac{1}{\sqrt{2}} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix}$ ,  $S = \begin{bmatrix} 1 & 0 \\ 0 & i \end{bmatrix}$ ,  $CNOT = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$

T gate:  $T = \begin{bmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{bmatrix}$

[Gottesman-Knill 98]:  $U$  Clifford  $\Rightarrow$  Yes.

$$|H\rangle = e_0 + e^{i\pi/4} e_1$$

[Bravyi-Smith-Smolin 16, Bravyi-Gosset 16]: If  $U$  consists of  $m$  T gates, then circuit can be strongly simulated with cost quadratic in  $\chi(|H\rangle^{\otimes m})$ .

... and can be weakly simulated with cost quadratic in  $\chi_\delta(|H\rangle^{\otimes m})$ .

# Chapter 5: New techniques for bounding stabilizer rank

$\chi(|H\rangle^{\otimes m})$

- [Huang-Newman-Szegedy 20]:  $\chi(|H\rangle^{\otimes m})$  super-polynomial unless P=NP.
- [Bravyi-Smith-Smolín 16]:  $\chi(|H\rangle^{\otimes m}) \geq \Omega(\sqrt{m})$ .
- [Peleg-Shpilka-Volk 21]:  $\chi(|H\rangle^{\otimes m}) \geq \Omega(m)$ .
- [Qassim-Pashayan-Gosset 21]:  $\chi(|H\rangle^{\otimes m}) \leq 2^{\alpha m}$ , where  $\alpha = \frac{1}{4} \log_2(3)$ .

$\chi_\delta(|H\rangle^{\otimes m})$

[Lovitz-Steffan 21]: Alternate (simplified) proofs up to log factor

- [Peleg-Shpilka-Volk 21]:

There exists  $\delta > 0$  such that

$$\chi_\delta(|H\rangle^{\otimes m}) \geq \Omega(\sqrt{m}/\log m)$$

- [Bravyi-Gosset 16]:

$$\chi_\delta(|H\rangle^{\otimes m}) \leq O\left(\frac{1}{\delta^2} 2^{\alpha m}\right), \text{ where } \alpha \approx 0.228.$$



# Summary



- Algebraic statistics: View a **probability vector** for random variables  $X, Y, Z$  as a tensor  $T$ .

Chapter 7 • Splitting theorem  $\Rightarrow$  Generalization of Kruskal's theorem.

- Algebraic complexity theory: View a **bilinear map**  $X \times Y \rightarrow Z^*$ , e.g. matrix multiplication, as a tensor  $T$ .

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Chapter 6 • Decomposable correlation matrices

Chapter 3 • Generic local state discrimination with pre-shared entanglement

Chapter 5 • New techniques for bounding stabilizer rank